# A CLASS OF SINGLE SAMPLING PLANS BASED ON FUZZY OPTIMISATION 

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#### Abstract

The problem of determining single sampling lot tolerance per cent defective (LTPD) inspection plans is considered under the fuzzy environment of satisfying the consumer's risk closely. A solution procedure, under Poisson conditions, for the above non-symmetrical fuzzy mathematical programming model involving minimisation of a crisp objective function with fuzzy solution space is developed. Numerical examples and a sensitivity analysis are included.


## 1. Introduction

We consider a producer's final inspection of a series of lots under the industrial production process where each lot retains its identity such as lots of electronic equipment for a large computer or a missile. The producer wants to minimize the average total inspection for product of process average quality and at the same time wants to be reasonably sure that lots of bad quality are not marketed.

In designing an attribute sampling plan for acceptance inspection, it is assumed that the producer knows his process average quality level $p_{1}$ under normal manufacturing conditions and that he occasionally produces lots of bad quality. The decision maker (DM) may then select lot tolerance fraction defective, $p_{2}$ (also called LTPD quality level), say $p_{2}>p_{1}$ and a risk $P\left(p_{2}\right)=\beta$ of accepting the lots of this quality where $P(p)$ is the operating characteristic ( OC ) of the acceptance sampling plan.

In industrial applications, a single sampling plan (SSP) is generally employed. A SSP is the following: From each lot of size $N$, a random sample of size $n$ is taken. If the number of defectives in the sample is less than or equal to the acceptance number $c$, the lot is accepted; otherwise the lot is rejected.

We consider the class of SSPs for which rejected lots are 100 per cent inspected. In the above class the classical model of Dodge Romig [6] LTPD SSP is to find the sample size $n$ and the acceptance number $c$ satisfying the following criteria:

$$
\begin{equation*}
\text { Minimize } I\left(p_{1}, n, c\right)=n+(N-n)\left(1-P\left(p_{1}\right)\right) \tag{1}
\end{equation*}
$$

subject to $P\left(p_{2}, n, c\right) \leqslant \beta$,
and $\quad n, c \geqslant 0$, integer,
where $I\left(p_{1}, n, c\right)$ is the average total inspection (ATI).
We note that the above optimisation problem (1) through (3) is a nonlinear integer programming (NLIP) problem and may be viewed as a nonlinear integer knapsack problem (see Garfinkel and Nemhauser [7]). Since it is a NLIP problem, we have replaced the equality sign in (2), which does not change the solution. For applications of mathematical programming techniques in designing SSPs, see Chakraborty [2,4,5].

In industrial applications of sampling inspection plans, one faces imprecisions, for example, in the specification of LTPD quality, $p_{2}$, which is termed rejectable or unsatisfactory quality, and in the requirement of consumer's risk to be satisfied as close as possible to $\beta$ (see Chakraborty [3] and Hald [8]). Now all imprecisions cannot be treated by randomnes: and modelled by using probability theory. From the above, we see that the imprecisions are often fuzzy in character and require to be modelled using fuzzy set theory. When the environment is really fuzzy in nature. modelling the problem with erroneous assumption of randomness produces sampling plans with appreciably larger ATI.

We consider the Dodge-Romig SSPs and assume that the DM wants a sampling plan with consumer's risk is to be satisfied closely. We model the problem as a Fuzzy Mathematical Programming (FMP) and derive a method for obtaining the solution. We may note that the problem is nonsymmetrical (see Zimmermann [10,11]) since the objective function (1) is crisp and the constraint (2) is a fuzzy set.

We shall restrict our discussions for the industrial applications where $p_{2} \leqslant 0.1, p_{1} / p_{2}<0.5$ and $n / N \leqslant 0.1$ such that Poisson distribution will provide sufficiently accurate approximation for $P\left(p_{1}\right)\left(P\left(p_{2}\right)\right)$ for which binomial (hypergeometric) distribution is involved. So $P(p)=G(c, n p)=$ $\sum_{x=0}^{c} e^{-n p}(n p)^{x} / x!$. For extension to the appropriate distributions involved,
for notations and explanation of the terms, see Hald [8] and Chakraborty $[2,3]$.

## 2. Fuzzy Mathematical Programming Model

The Dodge Romig [6] SSP when the DM wants to minimize the ATI at the process average $p_{1}$ subject to satisfying the consumer's risk closely around $\beta$ is to find non-fuzzy non-negative integer pair ( $n, c$ ) which minimizes

$$
\begin{equation*}
I\left(p_{1}, n, c\right)-n+(N-n)\left(1-G\left(c, n p_{1}\right)\right) \tag{4}
\end{equation*}
$$

subject to $G\left(c, n p_{2}\right) \lesssim \beta$.

$$
\begin{equation*}
n, c \geqslant 0, \text { integer. } \tag{5}
\end{equation*}
$$

The symbol $\underset{\sim}{\sim}$ refers to fuzzified version of $\leqslant$ sign and means approximately less than or equal to or essentially smaller than or closely around stated value.

Clearly the model is a nonlinear integer fuzzy mathematical programming (NLIFMP) problem.For details of fuzzy set theory and FMP problem, see Zimmermann [9, 10] and references in Chakraborty [3].

An optimum value of a crisp function over a crisp domain attains at a precise point, called an optimum decision. The decision models involving fuzziness, the optimum decision according to Bellman and Zadeh [1, p. 150], is often considered to be a crisp set which contains those elements of fuzzy set decision attaining the maximum degree of membership. In the model (4) through (6), the objective function (4) is crisp which is to be minimised and it induces an order of the decision alternatives. The constraint (5) defines a fuzzy decision space. It is seen that the assumption of Bellman and Zadeh [1] about the symmetry of objective function and constraint sets are not valid in this model. This is a non-symmetrical case and following Zimmermann [10, 11], we will transform the model to a symmetrical case so that we may solve the problem by adapting the method given in Chakraborty [3].

## Membership functions

Following Zimmermann [10, p. 231, 11, p. 104], we define:
Let $R$ be a fuzzy feasible region, $S(R)$ support of $R$ and $R_{1}$ a-level cut of $R$ for $a=1$. The membership function of objective function (4) given solution space $R$ is defined as

$$
\mu_{1}(n, c)=\left\{\begin{array}{cl}
1, & \text { if } I\left(p_{1}, n, c\right) \leqslant \inf _{S(R)} I  \tag{7}\\
\inf I-I\left(p_{1}, n, c\right) & \text { if inf } I<I\left(p_{1}, n . c\right)<\inf I, \\
\frac{R_{1}}{\inf I-\inf I}, & S(R) \\
R_{1}(S R) & \text { if inf } I<I\left(p_{1}, n, \mathrm{c}\right) \\
0 . & R_{1}
\end{array}\right\}
$$

Assuming that the $D M$ specifies the upper tolerance limit of $\beta$ as $\beta_{C}$. following Chakraborty [3] we define the fuzzy set corresponding to fuzzy constraint (5) by the membership function

$$
\mu_{1}=\left\{\begin{array}{cl}
1, & \text { if } P\left(p_{2}, n, c\right) \leqslant \beta  \tag{8}\\
\frac{\beta_{U}-P\left(p_{2}, n, c\right)}{\beta_{U}-\beta}, & \text { if } \beta<P\left(p_{2}, n, c\right)<\beta_{U} \\
0, & \text { if } \beta_{U} \leqslant P\left(p_{2}, n, c\right)
\end{array}\right\}
$$

Now symmetry is achieved between the objective function (4) and the constraint (5) and following Chakraborty [3] with minimum operator to aggregate the membership functions of fuzzy sets, we obtain the following model for the Dodge Romig problem:
$\left.\underset{n, c \geqslant 0, \operatorname{int}}{\operatorname{Maximise}[\min }\left\{\begin{array}{l}\inf I-I\left(p_{1}, n, c\right) \\ \frac{R_{1}}{\inf I-\inf I}=\lambda_{1},-\frac{\beta_{U}-P\left(p_{2}, n, c\right)}{\beta_{U}-\beta}=\lambda_{2} \\ R_{1} \quad S(R)\end{array}\right\}=\lambda\right]$.
subject to

$$
\begin{gather*}
\inf _{S(R)} \leqslant I\left(p_{1}, n, c\right)<\inf _{R_{1}} I  \tag{1}\\
\beta \leqslant P\left(p_{2}, n, c\right) \leqslant \beta_{U},  \tag{11}\\
0<\lambda<1 \tag{12}
\end{gather*}
$$

$$
\begin{equation*}
n, c>0, \text { integer. } \tag{13}
\end{equation*}
$$

This problem can be rewritten in the equivalent optimisation problem: find $n, c, \lambda$, which maximises $\lambda$, subject to

$$
\begin{equation*}
\lambda<\frac{\inf _{R_{1}} I-1\left(p_{1}, n, c\right)}{\inf _{R_{2}} I-\inf I}, \tag{14}
\end{equation*}
$$

$$
\begin{equation*}
\lambda<\frac{\beta_{U}-P\left(p_{2}, n, c\right)}{\beta_{U}-\beta} \tag{15}
\end{equation*}
$$

and inequalities (10) through (13).
Let $\quad I_{1} \equiv \inf _{R_{1}} I \quad$ and $\quad I_{o} \equiv \inf _{S(R)} I$.
Then the above optimisation problem (14), (15), (16), (10), (11), (12), (13) can be expressed as:

Find $n . c, \lambda$, which maximises $\lambda$, subject to

$$
\begin{align*}
& \lambda\left(I_{1}-I_{0}\right)+I\left(p_{1}, n, c\right)<I_{1}  \tag{16}\\
& \lambda\left(\beta_{I J}-\beta\right)+P\left(p_{2}, n, c\right)<\beta_{U}  \tag{17}\\
& \quad I_{0}<I\left(p_{1}, n, c\right)<I_{1} \tag{18}
\end{align*}
$$

and inequalities (11), (12) and (13).
This is a non-fuzzy mixed integer nonlinear programming (NFMINLP) problem. Noting that the values of $I_{1}$ and $I_{0}$ can be found by solving two non-fuzzy optimisation problems. the above NFMINLP problem can be solved by standard methods.

## 3. Solution Procedure and Examples

It is obvious that the values of $I_{0}$ and $I_{1}$ can be obtained by solving the following two NLIP problems respectively, viz.,
$\operatorname{Min} I\left(p_{1}, n, c\right)=n+(N-n)\left(1-\boldsymbol{G}\left(c, n p_{1}\right)\right)$,
subject to $\quad G\left(c, n p_{2}\right)<\beta$,

$$
\begin{equation*}
n, c \geqslant 0, \text { integer } \tag{20}
\end{equation*}
$$

and $\operatorname{Min} I_{1}\left(p_{1}, n, c\right)=n+(N-n)\left(1-G\left(c, n p_{1}\right)\right)$,
subject to $\beta \leqslant G\left(c, n p_{2}\right) \leqslant \beta_{U}$,

$$
\begin{equation*}
n, c>0, \text { integer. } \tag{23}
\end{equation*}
$$

Note that for a given $c, I(p, n, c)$ is an increasing function of $n$, since for a fixed $c, G(c, n p)$ is a decreasing function of $n$. Thus for a fixed $c$, the minimum $I(p, n, c)$ will be for the minimum feasible $n$. Clearly the optimum $n$ for the problem (19) through (21) for a given $c$ is the smallest $n$ satisfying
$G\left(c, n p_{2}\right) \leqslant \beta$. Similarly the optimum $n$ for the problem (22) through (24) for a given $c$ is the smallest $n$ satisfying $G\left(c, n p_{2}\right)<\beta_{l}$.

Let $n p=m$ and $m_{\beta}(c)$ be the $\beta$-fractile of the Poisson OC, $G(c, m)$. The values of $m_{\beta}(c)$ are given in Table 1 of Hald [8] for $c=0(1) 39$ for 11 values of $\beta$ and in Table 1 of Chakraborty [3] for $c=0(1) 9$ for 28 values of $\beta$. It is easily seen that the minimum $n$ satisfying $G(c, n p)<\beta$ i, given by

$$
\begin{equation*}
n(c)=1-m_{\beta}(c) / p \tag{25}
\end{equation*}
$$

where for a real number $y,\lceil y$ is the smallest integer $>1$.
For each $c$, we can find $n(c)$ and from the nature of $I\left(p_{1}, n, c\right)$, the optimum solution for the problems (19) through (21) as well as (22) through (24) can be obtained by direct search method; see Hald [8, p. 101].

Let the optimal solution for the problem (19) through (21) be ( $n_{0}, c_{o}$ ) with $I\left(p_{1}, n^{o}, c^{o}\right)_{o p t}=I_{1}$ and that for the problem (22) through (24) be ( $n_{0}, c_{0}$ ) with $I\left(p_{1}, n_{o}, c_{o}\right)_{\text {opt }}=I_{0}$.

Finally we consider the problem (16) through (18) together with (11) through (13). The solution procedure is similar with the difference that the set of feasible $n$ 's for a given $c$, is an interval instead of a singleton. For an arbitrary $c$, say $c_{o}$, we obtain the set of feasible $n$ 's satisfying (11)by using Table of $m_{\beta}(c)$. For each $n$ of the set, we calculate $\lambda=\lambda_{1}$ for (16) and $\lambda=\lambda_{2}$ for (17) and obtain the minimum $\lambda$ of the two, i.e., $\lambda=\min , \lambda_{1}, \lambda_{2}$. The $n$ which gives maximum of the minimum $\lambda$ 's is the optimum $n$ for this $c=c_{o}$. The procedure is repeated for $c$ 's near $\left\{c_{o}, c_{0}^{o}\right\}$. and the SSP giving highest value of $\lambda$ is the optimum $\operatorname{SSP}\left(n^{*}, c^{*}\right)$.

The search can be curtailed by noting, say, $\lambda_{o}$ be the maximum $\lambda$ corres ponding to $c_{o}$ and updating inequalities (16), (17) and (12). i.c.,

$$
\begin{align*}
& I\left(p_{1}, n, c\right) \leqslant I_{o}-\lambda_{o}\left(I_{1}-I_{o}\right)=I_{o}^{\prime}  \tag{26}\\
& P\left(p_{2}, n, c\right) \leqslant \beta_{U}-\lambda_{o}\left(\beta-\beta_{U}\right)=\beta_{U}^{\prime}  \tag{27}\\
& \lambda_{o} \leqslant \lambda \leqslant 1 \tag{28}
\end{align*}
$$

This will result in only a small number of feasible $c$ 's. Also the cardinality of the set of feasible $n$ 's for each feasible $c$ will be very small.

## Numerical Examples

Example 1. We consider the example given in Hald [8, p. 101]: $N=2,000$.
$p_{1}=0.02, p_{2}=0.10, \beta=0.10$. In addition, let $\beta_{U}=0.15$. The problem is to find the optimum LTPD $\operatorname{SSP}\left(n^{*}, c^{*}\right)$.

Solution. Solving the optimisation problems, we obtain $\left(n^{\circ}, c^{\circ}\right)=$ $(93,5)$ with $I_{1}=115.93$ and $\left(n_{o}, c_{o}\right)=(85,5)$ with $I_{o}=100.32$. For $c_{o}=5$, we obtain the set of feasible sample size $n$ 's satisfying (11) as $\left[n_{1}, n_{2}\right]=$ [85,92] and obtain the optimum $n$ for this $c=c_{0}$ as $n=98$ with $\lambda_{o}=$ 0.517 by constructing Table 1. Using $\lambda_{0}=0.517$, we find from (26), (27) and (28), $I\left(p_{1}, n, c\right)<107.86, P\left(p_{2}, n, c\right)<0.124$ and $0.517<\lambda \leqslant 1$.

For $c=4$, sample size $n$ 's satisfying (11) is $\left[n_{1}, n_{2}\right]=[73,79]$.
For $c=4$, the set of $n$ 's satisfying (11) and (27) is [77, 79].
However, for $c=4$, the set of $n$ 's satisfying (11), (26) and (27) is empty. Similarly for $c=6$, the set of sample size $n$ 's satisfying (11) is [98, 105]; satisfying (11) and (29) is [101, 105]. However, for $c=6$, the set of feasible sample size is empty. Clearly the optimum LTPD SSP is $n^{*}=89, c^{*}=5$ with $I^{*}\left(p_{1}, n^{*}, c^{*}\right)=107.86$ with $\lambda^{*}=0.517$. Note also that $P\left(p_{2}, n^{*}, c^{*}\right)$ $=0.1219$

In this particular example, the DM takes an additional risk of $2 \%$ for a lot being rejected against a saving of $7 \%$ inspection effort per lot.

Example 2: Same problem as in Example 1 except $N=5,000$.
TABLE 1-OPTIMAL SAMPLE SIZE FOR $c_{o}=5$

| $n$ | $P\left(p_{2}\right)$ | $I\left(p_{1}\right) \quad \lambda_{1}=\frac{I_{1}-I\left(p_{1}\right)}{I-I_{o}}$ | $\lambda_{2}=\frac{\beta_{U}-P\left(p_{2}\right)}{\beta_{U}-\beta}$ | $\lambda=\min \left\{\lambda_{1}, \lambda_{2}\right\}$ |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 85 | .1496 | 100.32 | 1.000 | .008 | .008 |
| 88 | .1284 | 105.92 | 0.641 | .432 | .432 |
| 89 | .1219 | 107.86 | 0.517 | .562 | .517 |
| 90 | .1157 | 109.82 | .0391 | .685 | .391 |
| 92 | .1041 | 113.86 | 0.133 | .918 | .133 |

TABLE 2 ---OPTIMUM SAMPLE SIZE FOR DIFFERENT $c$


Solution. Here $\left(n^{\circ}, c^{\circ}\right)=(118,7)$ with $I_{1}=132.75$ and

$$
\left(n_{0}, c_{o}\right)=(98,6) \text { with } I_{0}=117.96
$$

Proceeding as in Example 1, we search for plans with $c$ 's near $\{6,7\}$ and check for each feasible $n$ corresponding to each $c$. From Table 2, the optimum LTPD SSP is $n^{*}=101, c^{*}=6, I^{*}=124.41$ and $\lambda^{*}=0.520$.

## 4. Sensitivity Analysis

## Effect of Lot Size

Because of discreteness of $c$, the plan ( $n^{*}, c^{*}$ ) which is optimum for $N$ will also be optimum for some neighbouring values of $N$. The optimum SSP $(89,5)$ for the problem of Example 1, is optimum for all $N \in[1399$. 2436]. For a given set of parameters ( $p_{1}, p_{2}, \beta$ and $\beta_{U}$ ), it is seen that the optimum LTPD SSPs have the following properties: (i) $c$ and $n$ are increasing functions of $N$, (ii) ATI is increasing with $N$. A typical example is given in Table 3.

## Effect of Tolerance Limit of Consumer's Risk

For a given set of parameters $\left(N, p_{1}, p_{2}, \beta\right)$, as the tolerance limit $\beta_{t}$ increases within reasonable limit, the optimum decision number remains same and the optimum sample size and ATI decreases slowly. However if $\beta_{U}$ is increased appreciably, the optimum decision number may also decrease by one unit. A typical example is given in Table 4.

TABLE 3-EFFECT OF LOT SIZE ON THE FUZZY LTPD SSPs
$\left(p_{1}=0.02, p_{2}=0.10, \beta=0.10, \beta_{U}=0.15\right)$

| $N$ | $c^{*}$ | $n^{*}$ | $P\left(p_{2}\right)$ | $I\left(p_{1}\right)$ | $\lambda^{*}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 500 | 3 | 64 | 0.1189 | 81.93 | 0.508 |
| 1000 | 4 | 76 | 0.1249 | 94.05 | 0.502 |
| 1399 | 5 | 89 | 0.1219 | 101.93 | 0.502 |
| 1500 | 5 | 89 | 0.1219 | 102.92 | 0.515 |
| 1800 | 5 | 89 | 0.1219 | 105.88 | 0.516 |
| 2000 | 5 | 89 | 0.1219 | 107.86 | 0.517 |
| 2200 | 5 | 89 | 0.1219 | 109.83 | 0.518 |
| 2436 | 5 | 89 | 0.1219 | 112.16 | 0.495 |
| 5000 | 6 | 101 | 0.1240 | 124.41 | 0.520 |
| 10000 | 7 | 113 | 0.1249 | 136.02 | 0.502 |

TABLE 4-EFFECT OF TOLERANCE LIMIT ON THE FUZZY LTPD SSPs

$$
\left(N=2,000, \quad p_{1}=0.02, \quad p_{2}=0.10, \quad \beta=0.10\right)
$$

| $\beta_{U}$ | $c^{*}$ | $n^{*}$ | $P\left(p_{2}\right)$ | $\mathrm{I}\left(p_{1}\right)$ | $\lambda^{*}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0.11 | 5 | 92 | 0.1041 | 113.86 | 0.504 |
| 0.13 | 5 | 90 | 0.1157 | 109.82 | 0.477 |
| 0.15 | 5 | 89 | 0.1219 | 107.86 | 0.517 |
| 0.18 | 5 | 87 | 0.1352 | 104.02 | 0.560 |
| 0.20 | 5 | 86 | 0.1422 | 102.15 | 0.566 |

## 5. Concluding Remarks

In the classical Dodge Romig [6] model, one of the objectives of the DM of minimum consumer's risk is achieved by fixing it at a specified value of $\beta$ and then the second objective of minimum average amount of inspection is achieved by minimizing this objective function subject to the above constraint. In the present model, the DM specifies an interval $\left[\beta, \beta_{U}\right]$ for the risk that may be taken for allowing the lot to be marketed with a view
to reducing the amount of inspection per lot. This non-symmetric fuzzy model is able to scale the objective function with the constraint so that a compromise solution is obtained and thereby reduces the amount of inspection in comparison with the non-fuzzy model.

The model can be easily extended to the case when the DM can specify the weight or cost per unit risk per lot and that for reduced inspection per unit per lot. The DM may also consider addition operator to aggregate the fuzzy sets and design the optimum SSP by maximising the weighted fuzzy achievement function (see Chakraborty [3]).

## Acknowledgements

The author wishes to express his apprecitation to Professor A.C. Mukhopadhyay for generously giving many valuable suggestions. Thank, are due to the referees for valuable suggestions that resulted in substantial improvement of the paper.

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