# COMPUTER AIDED-CONSTRUCTION OF D-OPTIMAL $2^{m}$ FRACTIONAL FACTORIAL DESIGNS OF RESOLUTION $V$ 

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## Summary

A new exchange algorithm for construction of $2^{m} D$-optimal fractional factorial design (FFD) is devised. This exchange algorithm is a modification of the one due to Fedorov $(1969,1972)$ and is an improvement over similar algorithm due to Mitchell (1974) and Galil \& Kiefer (1980). This exchange algorithm is then used to construct $54 D$-optimal $2^{m}$-FFD's of resolution $V$ for $m=4,5,6$.

Key words: Fractional factorial design; D-optimality; A-optimality; exchange algorithm.

## 1. Introduction

A fractional factorial design (FFD) is said to be of resolution $V$ if it permits estimation of the mean, all main effects and all two factor interactions, under the assumption that all interactions between three or more factors are negligible in magnitude. Thus, a total number of parameters to be estimated in a resolution $V$ design for $2^{m}$-FFD is $p=1+m+m C_{2}$.

An orthogonal resolution $V$ plan for a $2^{m}-\mathrm{FFD}$ is equivalent to an orthogonal array $\mathrm{OA}(n, m, 2,4)$, i.e. an OA with $n$ assemblies, $m$ constraints, 2 symbols and strength 4 (cf. Rao (1947)). When such an OA is available, it provides an optimal resolution $V$ design with respect to

[^0]$D$-, $A$-, $E$-optimality criteria. In fact, this design is universally optimal (Kiefer, 1980). However, an OA providing an optimal resolution $V$ design exists only if $n=0(\bmod 4)$, a condition which may not always be possible to satisfy in practice. Thus, the need for optimal resolution $V$ design for other values of $n$ arises.

Srivastava \& Chopra (1971) have considered $A$-optimal resolution $V$ designs for $2^{m}$-FFD for $m=4,5$ and 6 for practical values of $n$ in the class of balanced designs. These designs are balanced in the sense that the variance-covariance matrix $V$ of the parameter estimates is invariant under the permutation of the $m$ factors. However, the balanced designs form only a subclass of designs and one may like to study the optimality of designs in the entire class. Such a study has been partially made by Kuwada (1982) who constructed optimal resolution $V$ design $2^{m}$-FFD for $m=4,5$ and 6 with respect to $A$-optimal criterion. Some of these designs are in fact superior to the corresponding designs of Srivastava \& Chopra who restricted their attention to balanced designs.

The purpose of this study is to devise a computer algorithm for the construction of $D$-optimal resolution $V 2^{m}$-FFD. This algorithm is then used actually to construct $54 D$-optimal $2^{m}$-FFD of resolution $V$ for $m=4$, 5 and 6.

## 2. Some Matrix Results

In this section, we give some results in matrix algebra, which will be used in the sequel.

Consider the usual full rank linear model $y=X \beta+e$ where $y$ is an $n$-component column vector of observations, $X$ is an $n \times p$ matrix of known elements, $\beta$ is a $p \times 1$ vector of unknown parameters and $e$ is a $p \times 1$ vector of random residual components with $E(e)=0$ and $D(e)=\sigma^{2} I$ where $E$ stands for the expectation operator and $D$ denotes the dispersion matrix. The $n$ rows of $X$ are $n$-dimensional vectors $x_{i}^{\prime}, i=1,2, \ldots, n$. These $n$ vectors can be considered as $n$ points in $R^{p}$. Let $M=X^{\prime} X$ and assume $M$ to be non-singular.

If $x^{\prime}$ is a row vector to be augmented to $X$, we have:

$$
\begin{align*}
\operatorname{det}\left(M+x x^{\prime}\right) & =\operatorname{det}(M)\left(1+x^{\prime} M^{-1} x\right)  \tag{2.1}\\
\left(M+x x^{\prime}\right)^{-1} & =M^{-1}+w u u^{\prime} \tag{2.2}
\end{align*}
$$

where $w=-\left(1+x^{\prime} M^{-1} x\right)^{-1}$ and $u=M^{-1} x$.

Now let $M_{x}=M+x x^{\prime}$. If $x_{i}^{\prime}$ is a row vector to be removed from the current $X$, we have:

$$
\begin{align*}
\operatorname{det}\left(M_{x}-x_{i} x_{i}^{\prime}\right) & =\operatorname{det}\left(M_{x}\right)\left(1-x_{i}^{\prime} M_{x}^{-1} x_{i}\right)  \tag{2.3}\\
\left(M_{x}-x_{i} x_{i}^{\prime}\right)^{-1} & =M_{x}^{-1}+w_{i} u_{i} u_{i}^{\prime} \tag{2.4}
\end{align*}
$$

where $w_{i}=\left(1-x_{i}^{\prime} M_{x}^{-1} x_{i}\right)^{-1}$ and $u_{i}=M_{x}^{-1} x_{i}$.
Now, let $x^{\prime}$ be a row vector augmented to $X$ and $x_{i}^{\prime}$ be a row vector removed from $X$ simultaneously, i.e. $x_{i}^{\prime}$ and $x^{\prime}$ are exchanged. Then we have:

$$
\begin{equation*}
\operatorname{det}\left(M+x x^{\prime}-x_{i} x_{i}^{\prime}\right)=\operatorname{det}(M)\left\{1+D\left(x_{i}, x\right)\right\} \tag{2.5}
\end{equation*}
$$

where:

$$
\begin{equation*}
D\left(x_{i}, x\right)=x^{\prime} M^{-1} x-x_{i}^{\prime} M^{-1} x_{i}\left(1+x^{\prime} M^{-1} x\right)+\left(x^{\prime} M^{-1} x_{i}\right)^{2} . \tag{2.6}
\end{equation*}
$$

## 3. Method of Construction

The model for $2^{m}$-FFD of resolution $V$ is the usual full rank linear model $y=X \beta+e$ as in Section 2. The $i$ th row of design matrix $X$ is a $p$-dimensional row vector $x_{i}^{\prime}$ :

$$
x_{i}^{\prime}=\left(1, x_{1 i}, x_{2 i}, \ldots, x_{m i}, x_{1 i} x_{2 i}, \ldots, x_{m-1 i} x_{m i}\right)
$$

where $x_{h i}= \pm 1, h=1,2, \ldots, m$ and $i=1,2, \ldots, n$.
The total number of candidate vectors $x$ is $2^{m}$. Our problem is that for a given $n$, we have to choose $n$ vectors $x$ 's out of $2^{m}$ candidate vectors such that $\operatorname{det}\left(X^{\prime} X\right)$ is maximized. Here, $n$ is not necessarily $\leq 2^{m}$ and the $x$ 's are not necessarily distinct.

Let $M=X^{\prime} X$. The proposed exchange algorithm (EA) for finding $D$-optimal $2^{m}$-FFD of resolution $V$ consists of the following steps:
(i) Start with a randomly chosen non-singular $n$-point design. Compute $M, M^{-1}$ and $\operatorname{det}(M)$.
(ii) Find a vector $x$ among $2^{m}$ candidate vectors such that $x^{\prime} M^{-1} x$ is maximum. This $x^{\prime} M^{-1} x$ is $V_{\max } / \sigma^{2}$, where $V_{\max }$ is the maximum variance of the predicted response of the current $n$-point design.
(iii) Find a vector $x_{i}$ among $n$ vectors of the current $n$-point design such that $D\left(x_{i}, x\right)$ is maximum. $D\left(x_{i}, x\right)$ is calculated by (2.6).
(iv) If $D\left(x_{i}, x\right)$ is less than a chosen positive small number say $10^{-5}$, then terminate. Otherwise exchange vector $x_{i}$ with $x$. Update $\operatorname{det}(M)$ by (2.5) and $M^{-1}$ by (2.2) and (2.4). Then return to step (ii).

TABLE 1
D-optimal $2^{4}$-FFD of resolution $V$

| $n$ | $\left\|X^{\prime} X\right\|$ | $V_{\text {mas }}$ | $\operatorname{tr} V$ | $\operatorname{tr} V_{k}$ | $\operatorname{tr} V_{t}$ |
| :--- | :---: | :---: | :---: | :---: | :---: |
| 11 | $3.86547 \mathrm{E}+10$ | 2.55556 | 1.48611 | - | 1.4861 |
| 12 | $1.37439 \mathrm{E}+11$ | 2.50000 | 1.31250 | 1.31250 | 1.3125 |
| 13 | $4.81036 \mathrm{E}+11$ | 2.42857 | 1.14286 | 1.14286 | 1.2639 |
| 14 | $1.64927 \mathrm{E}+12$ | 2.33333 | 0.97917 | 0.97917 | 1.1875 |
| 15 | $5.49756 \mathrm{E}+12$ | 2.20000 | 0.82500 | 0.82500 | 0.8250 |
| 16 | $1.75922 \mathrm{E}+13$ | 0.68750 | 0.68750 | 0.68750 | 0.6875 |
| 17 | $2.96868 \mathrm{E}+13$ | 0.68519 | 0.66204 | 0.66204 | 0.6620 |
| 18 | $5.00278 \mathrm{E}+13$ | 0.68269 | 0.63668 | 0.63668 | 0.6375 |
| 19 | $8.41814 \mathrm{E}+13$ | 0.68000 | 0.61143 | 0.61143 | 0.6270 |
| 20 | $1.41425 \mathrm{E}+14$ | 0.67708 | 0.58631 | 0.58631 | 0.5863 |
| 21 | $2.37181 \mathrm{E}+14$ | 0.63975 | 0.56134 | 0.63908 | 0.5613 |
| $22^{*}$ | $3.89639 \mathrm{E}+14$ | 0.65714 | 0.53780 | 0.63720 | 0.5384 |
| 23 | $6.45688 \mathrm{E}+14$ | 0.65517 | 0.51365 | 0.63575 | 0.5136 |
| 24 | $1.06873 \mathrm{E}+15$ | 0.58333 | 0.48958 | 0.63462 | 0.4896 |
| 25 | $1.69215 \mathrm{E}+15$ | 0.57895 | 0.46930 | 0.63370 | 0.4693 |
| $26^{*}$ | $2.68006 \mathrm{E}+15$ | 0.60256 | 0.44888 | 0.63294 | 0.4518 |
| 27 | $4.29497 \mathrm{E}+15$ | 0.53600 | 0.42750 | 0.63230 | 0.4275 |
| 28 | $6.59707 \mathrm{E}+15$ | 0.53333 | 0.41042 | - | 0.4104 |

* trace $V$ is strictly less than either of trace $V_{k}$ or trace $V_{s}$.

This EA, like Mitchell's DETMAX (1974) and Galil \& Kiefer's modified DETMAX or MD (1980) is another version of Fedorov's EA (1969, 1972) (cf. St. John \& Draper (1975)). One advantage of this EA over DETMAX and MD is that double precision is not required in the computation of $\operatorname{det}\left(M+x x^{\prime}-x_{i} x_{i}^{\prime}\right)$ since the straightforward formula (2.5) is used. In DETMAX, for example $\left(M+x x^{\prime}\right)^{-1}$ has to be evaluated before the evaluation of $\operatorname{det}\left(M+x x^{\prime}-x_{i} x_{i}^{\prime}\right)$. Another advantage of this EA over DETMAX and MD is that an array of length $2^{m}$ need not be maintained in the computer to store $2^{m}$ values of $x^{\prime} M^{-1} x$.

Like all previous EA's, this new EA does not always guarantee $D$ optimality as it may get "trapped" in the local optimum. In order to get a good design for given $m$ and $n$, several tries should be made, each try

TABLE 2
D-optimal $2^{5}$-FFD of resolution $V$

| $n$ | $\left\|X^{\prime} X\right\|$ | $V_{\text {max }}$ | $\operatorname{tr} V$ | $\operatorname{tr} V_{k}$ | $\operatorname{tr} V_{\boldsymbol{z}}$ |
| :--- | :---: | :---: | :---: | :---: | :---: |
| 16 | $1.84467 \mathrm{E}+19$ | 1.00000 | 1.00000 | 1.00000 | 1.0000 |
| 17 | $3.68935 \mathrm{E}+19$ | 1.00000 | 1.96875 | 0.96875 | 0.9687 |
| 18 | $7.37870 \mathrm{E}+19$ | 1.00000 | 1.93750 | 0.93750 | 0.9398 |
| 19 | $1.47574 \mathrm{E}+20$ | 1.00000 | 0.90625 | 0.90625 | 0.9296 |
| 20 | $2.95148 \mathrm{E}+20$ | 1.00000 | 0.87500 | 0.87500 | 0.9194 |
| 21 | $5.90296 \mathrm{E}+20$ | 1.00000 | 0.84375 | 0.84375 | 0.8437 |
| 22 | $1.18059 \mathrm{E}+21$ | 1.00000 | 0.81250 | 0.94643 | 0.8125 |
| $23^{*}$ | $2.36118 \mathrm{E}+21$ | 1.00000 | 0.78125 | 0.94531 | 0.7979 |
| $24^{*}$ | $4.72237 \mathrm{E}+21$ | 1.00000 | 0.75000 | 0.94444 | 0.7881 |
| $25^{*}$ | $9.44473 \mathrm{E}+21$ | 1.00000 | 0.71875 | 0.94375 | 0.7815 |
| 26 | $1.88895 \mathrm{E}+22$ | 1.00000 | 0.68750 | 0.94318 | 0.6875 |
| 27 | $3.77789 \mathrm{E}+22$ | 1.00000 | 0.65625 | 0.65625 | 0.6563 |
| 28 | $7.55579 \mathrm{E}+22$ | 1.00000 | 0.62500 | 0.62500 | 0.6300 |
| 29 | $1.51116 \mathrm{E}+23$ | 1.00000 | 0.59375 | 0.59375 | 0.6199 |
| 30 | $3.02231 \mathrm{E}+23$ | 1.00000 | 0.56250 | 0.56250 | 0.5830 |
| 31 | $6.04463 \mathrm{E}+23$ | 1.00000 | 0.53125 | 0.53125 | 0.5313 |
| 32 | $1.20893 \mathrm{E}+24$ | 0.50000 | 0.50000 | 0.50000 | 0.5000 |

* trace $V$ is strictly less than either of trace $V_{k}$ or trace $V_{s}$.
with a different starting design. In this study, 10 tries are made for each design with given $m$ and $n$.


## 4. Results and Discussion

The values of $\operatorname{det}\left(X^{\prime} X\right)$ of 54 constructed $D$-optimal $2^{m}$-FFD of resolution $V$ for $m=4,5$ and 6 together with trace $V$ where $V=\left(X^{\prime} X\right)^{-1}$, $V_{m a x}$, trace $V_{k}$ and trace $V_{s}$ are given in Tables 1, 2 and 3. $V_{k}$ and $V_{s}$ stand for the variance-covariance matrix of the designs obtained by Kuwada and by Srivastava \& Chopra. For these designs, it was found that trace $V$ is always less than or equal to the lesser of trace $V_{k}$ and trace $V_{s}$. All in all, there are 14 new designs having trace $V$ strictly smaller than either of trace $V_{k}$ or trace $V_{s}$. As expected, none of the obtained designs is balanced in the sense of Srivastava \& Chopra.

TABLE 3
D-optimal $\mathbf{2}^{6}$-FFD of resolution $V$

| $\boldsymbol{n}$ | $\left\|X^{\prime} X\right\|$ | $V_{\text {mas }}$ | tr $V$ | tr $V_{\boldsymbol{k}}$ | tr $V_{\mathbf{t}}$ |
| :--- | :---: | :---: | :---: | :---: | :---: |
| $22^{*}$ | $6.27415 \mathrm{E}+28$ | 1.34667 | 1.15167 | - | 1.6249 |
| 23* $^{*}$ | $1.47233 \mathrm{E}+29$ | 1.34259 | 1.11569 | - | 1.1241 |
| 24 $^{*}$ | $3.44908 \mathrm{E}+29$ | 1.33816 | 1.07974 | - | 1.1145 |
| 25 $^{*}$ | $8.06451 \mathrm{E}+29$ | 1.33333 | 0.04382 | - | 1.1100 |
| $26^{*}$ | $2.17607 \mathrm{E}+30$ | 1.35111 | 0.00278 | - | 1.1012 |
| 27 | $5.64036 \mathrm{E}+30$ | 1.70222 | 0.97542 | 1.36458 | 0.9754 |
| $28^{*}$ | $1.52415 \mathrm{E}+31$ | 1.70175 | 0.92544 | 1.00000 | 0.9487 |
| $29^{*}$ | $4.11788 \mathrm{E}+31$ | 1.70130 | 0.87541 | 0.91518 | 0.9371 |
| 30 | $1.21694 \mathrm{E}+32$ | 2.33333 | 0.83333 | 0.83333 | 0.9279 |
| 31 | $4.05648 \mathrm{E}+32$ | 2.20000 | 0.75625 | 0.75625 | 0.7562 |
| 32 | $1.29807 \mathrm{E}+33$ | 0.68750 | 0.68750 | 0.68750 | 0.6875 |
| 33 | $2.19050 \mathrm{E}+33$ | 0.68519 | 0.67477 | 0.67477 | 0.6747 |
| 34 | $3.69140 \mathrm{E}+33$ | 0.68304 | 0.66209 | 0.66209 | 0.6633 |
| 35 | $6.21276 \mathrm{E}+33$ | 0.68103 | 0.64945 | 0.64945 | 0.6582 |
| 36 | $1.04439 \mathrm{E}+34$ | 0.67917 | 0.63686 | 0.63686 | 0.6532 |
| 37 | $1.75370 \mathrm{E}+34$ | 0.67742 | 0.62430 | 0.62430 | 0.6245 |
| 38 | $3.17438 \mathrm{E}+34$ | 0.67511 | 0.60877 | 0.61178 | 0.6087 |
| $39^{*}$ | $5.31744 \mathrm{E}+34$ | 0.67326 | 0.59627 | 0.66163 | 0.5992 |
| $40^{*}$ | $8.89748 \mathrm{E}+34$ | 0.67129 | 0.58381 | 0.66106 | 0.5939 |

* trace $V$ is strictly less than either of trace $V_{\mathbf{k}}$ or trace $V_{\mathbf{s}}$.

For $m=5$ it takes about $\frac{1}{2}$ minutes per try on an IBM AT-compatible personal computer with an 80287 math coprocessor. For $m=6$ it takes about $2 \frac{1}{2}$ minutes per try and 10 tries may not be enough for a particular value of $n$. Out of 10 tries, the best design with respect to $D$-optimality criterion is chosen. However, for $m=6$ and for some values of $n$, it is not always true that the chosen designs have smaller trace and smaller $V_{\text {max }}$ than the rejected designs because the choice is based on $\boldsymbol{D}$-optimality criterion.

In concluding, we may remark that although we have presented results for $m=4,5$ and 6 only, the algorithm can be used for any values of $m$, for any resolution and for any factorial. Of course, for higher values of $m$ and
greater number of levels, the computer time requirement will be greater.
A PASCAL program listing of about 200 statements for constructing the designs in this paper can be obtained from the first author.

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