# A new split-and-merge clustering technique

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Received 5 August 1991

#### Abstract

Chaudhuri, D., B.B. Chaudhuri and C.A. Murthy, A new split-and-merge clustering technique, Pattern Recognition Letters 13 (1992) 399-409.

A new clustering algorithm is developed for efficient classification of data in  $\mathbb{R}^{q}$  when there exists no a priori information about the number of clusters. The algorithm is based on a split-and-merge technique. The type-I splitting is guided by density of data over strips at different directions around the centroid of the data. The type-II splitting is the usual *K*-means clustering algorithm (K = 2) and rechecked with the help of a merging technique. A theorem on the convergence of this algorithm is proved.

Keywords. Cluster analysis, split-and-merge, K-means clustering.

#### 1. Introduction

Clustering is a useful and important technique in image processing and pattern recognition [1,2,5, 7,9]. There exist two classes of clustering techniques, namely hierarchical and non-hierarchical techniques. Among non-hierarchical techniques K-means and ISODATA are popular. Of them, ISODATA is a split-and-merge technique of achieving a prespecified number of clusters. Among other split-and-merge techniques Wishart [4] as well as Liu and Tsai [3] may be mentioned. The method in this paper also falls in this category.

The main difference of the proposed method with those of the others is that here we try to split the clusters by noting the density at different directions by observing the data over strips. In order to overcome some defects of this approach, another splitting approach, that is, the simple 2-means algorithm under a certain restriction is used. Also, for merging, it is tested whether the data at the boundary of a cluster is very close to the data at the boundary of the other cluster. The splitting and merging techniques are described in Section 2. Section 3 describes the proposed algorithm with the corresponding flowchart and the theorem of convergence criterion. The results on synthetic and real data (remotely sensed imagery) are presented in Section 4.

#### 2. Splitting techniques

In type-I splitting, strips of finite width at different directions around the centroid of the data are considered. The data is split across the *sparsely populated strip*. For a two-dimensional data set we consider four directions at the center of the data as

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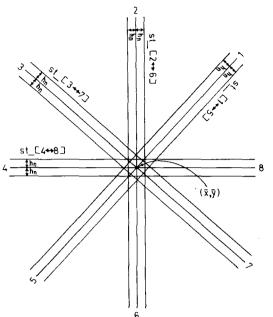


Figure 1. Strips in the four directions in a 2-dimensional data set.

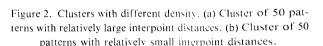
shown in Figure 1. They are named as  $1 \leftrightarrow 5$  (one diagonal),  $2 \leftrightarrow 6$  (vertical),  $3 \leftrightarrow 7$  (another diagonal), and  $4 \leftrightarrow 8$  (horizontal). The corresponding strips are denoted by  $St_{1} \rightarrow 5$ ,  $St_{2} \rightarrow 6$ , St\_ $[3 \leftrightarrow 7]$  and St\_ $[4 \leftrightarrow 8]$ , respectively. For a three-dimensional data set we consider  $2^{3-1} + 3 = 7$ directions at the centroid of the data  $(2^{3-1}=4)$ directions are 4 diagonals and 3 directions are 3 axes). The number of directions considered for a qdimensional data set is  $2^{q-1} + q$ . The strips are constructed along (i) the q axes and (ii) the diagonals (note that the number of diagonals in qdimensions is  $2^{q-1}$ ). But, observe that the greater the number of directions for strips the better the accuracy of the result is.

The width of a strip along any direction is to be found out before actually constructing the strip. The width is an important impediment in deciding whether the data is indeed sparsely populated in that direction. If the width is very large then all the points in the data set may belong to the strip. If it is very small then the strip may not be amenable for making any decision. In this connection a measure of finding the width is described below using an example.



(b)

June 1992



(a)

Example. Three data sets are shown in Figures 2(a), 2(b) and 3(a) where the number of points in these data sets are the same. Intuitively, the set of points in Figure 2(b) is a single cluster. The data in Figure 2(a) can be called a single cluster though interpoint distances are generally greater than in Figure 2(b). In Figure 3(a) it is intuitively clear that there are two clusters. Thus the sets of points in Figures 2(a) and 2(b) should not be split while those in Figure 3(a) are to be split. A way of splitting the set of points in Figure 3(a) at the center  $(X_0)$  is shown in Figure 3(b).

Draw a strip of width  $2h_n$  as shown in Figure 3(b). Count the number of points in the strip. If the number of points is less than some threshold, say  $\theta$ , then conclude that there are two clusters. Note that, if this procedure is to be followed in Figures 2(a) and 2(b) then  $h_n$  should be taken suitably.  $h_n$  should not be too big so that it could

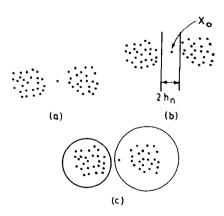


Figure 3. Cluster detection by splitting across the low-density strip. (a) Two clusters. (b) Lowest density strip of the data at the center  $X_0$ . (c) Splitted clusters.

include the entire set of points in the strip.  $h_n$  should not be too small so that the set of points in Figure 2(a) may be split. Note that  $h_n$  in a sense gives the connectivity of the set. In this regard a result [20] is stated in the Appendix. Based on this result we propose that the width of the strip should be  $2h_n$  where  $h_n = a\varepsilon_n$ ,  $\varepsilon_n = 1/n^p$ , 0 and a is a constant, n is the total number of points in the q-dimensional data set.

The type-I splitting is described below for  $\mathbb{R}^2$ . Let the given set of points be

$$S = \{(x_1, y_1), (x_2, y_2), \dots, (x_n, y_n)\} \subseteq \mathbb{R}^2.$$

Let

$$\bar{x} = \frac{1}{n} \sum_{i=1}^{n} x_i$$
 and  $\bar{y} = \frac{1}{n} \sum_{i=1}^{n} y_i$ .

Type-I splitting is applied at the mean point of the data set [i.e., at  $(\bar{x}, \bar{y})$ ] along a particular direction among the four directions discussed before.

Let  $b_1$ ,  $b_2$ ,  $b_3$  and  $b_4$  be the number of points in each of the strips  $St_{1} \leftrightarrow 5$ ,  $St_{2} \leftrightarrow 6$ ,  $St_{3} \leftrightarrow 7$  and  $St_{4} \leftrightarrow 8$ , respectively.

Let  $n_1 = \min\{b_1, b_2, b_3, b_4\}.$ 

Find the strip in which the number of points is  $n_1$ . If  $n_1$  is not very small compared to n, then there should not be any split in that direction. If

$$\frac{n_1}{n} \times 100 < l_1 \tag{1}$$

then split the set S along the corresponding direction. Inequality (1) is called the *splitting restriction*. Here  $l_1$  is a small quantity dependent on the width of the strip. If the number of points in a strip satisfies (1) then the strip is called a *sparsely populated strip*. Otherwise the strip is called a *densely populated strip*.

Note that if in two or more directions, the number of points,  $n_1$ , satisfies the *splitting restriction*, then one of the directions is chosen arbitrarily for splitting.

Type-I splitting can also be extended to three or more dimensions. Observe that for a two-dimensional data set, two straight lines are needed for constructing a strip (Figure 1). For higher-dimensional data sets, hyperdimensional strips need to be constructed. The number of hyperplanes needed for constructing a hyperdimensional strip in q dimensions is  $2^{q-1}$ . The inequality (1) would remain unchanged for *q*-dimensional data.

Note that *sparsely populated regions* need not always be present at the center of the data. They can occur at other places as well. In this context a splitting method, namely type-II splitting is incorporated in the algorithm.

The type-II splitting technique is the usual Kmeans method by Forgy [6] with some restriction imposed and K = 2. The selection of initial seed points is to be done suitably [14].

After the 2-means algorithm converges, let  $m_1$ and  $m_2$  be the mean points of the two subclusters  $S_1$  and  $S_2$  of S.

Let  $d_0 = ||m_1 - m_2||$ .

The almost equidistant point set, A, is the collection of those points whose differences of the distances from  $m_1$  and  $m_2$  are less than 10% of  $d_0$ , i.e.,

$$A = \left\{ z: \quad z \in S, \left| \left\| z - m_1 \right\| - \left\| z - m_2 \right\| \right| < \frac{d_0}{10} \right\}$$

and  $n_4 = #A$ . If

$$\frac{n_4}{n} \times 100 < l_2 \tag{2}$$

then  $S_1$  and  $S_2$  are said to be two different clusters. The inequality (2) is called the *almost equidistant restriction*. As before,  $l_2$  is a small quantity.

If  $(n_4/n) \times 100 \ge l_2$  then the merging criterion will have to be checked for deciding whether subclusters  $S_1$  and  $S_2$  remain divided or not.

#### Merging technique

Let  $S = \{z_1, z_2, ..., z_n\} \subseteq \mathbb{R}^d$ . Let  $S_1$  and  $S_2$  be such that  $S_1 \cap S_2 = 0$ ,  $S_1 \cup S_2 = S$  and they are the outcome of the 2-means algorithm on S. Let

$$m_i = \left(\sum_{z \in S_i} z\right) / \#S_i, \quad i = 1, 2.$$

Let

$$A = \left\{ z \colon z \in S, |\|z - m_1\| - \|z - m_2\| | < \frac{d_0}{10} \right\}$$

where  $d_0 = ||m_1 - m_2||$ . Let  $n_4 = #A$ . Let

$$m_3 = \frac{1}{n_4} \left( \sum_{z \in \mathcal{A}} z \right)$$

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and

$$H = \left\{ z : \quad \|m_3 - z\| \leq \frac{d_0}{5} \right\}.$$

The set *H* is called the *merging circle set* (Figure 4). Let

$$n_5 = \#H \cap S_1 \quad \text{and} \quad n_6 = \#H \cap S_2.$$

If

$$\frac{|n_5 - n_6|}{n} \times 100 < l_3 \tag{3}$$

then only  $S_1$  and  $S_2$  are to be merged.  $l_3$  must be very small. The inequality (3) is called the *merging* restriction.

It may be noted that

1. The above method basically verifies whether the points near the boundary of the two clusters have equal representation in both clusters (Figure 4).

2. The above method gives a criterion for merging  $S_1 \cup S_2$  which are products of the 2-mean algorithm on S.

3. If any one of  $S_1$  and  $S_2$ , say  $S_1$ , is further subdivided (say to  $S_{11}$  and  $S_{12}$ ), then also the same merging technique can be applied. The definitions  $m_i$ , A,  $d_0$ ,  $n_4$ ,  $m_3$ , H,  $n_5$  and  $n_6$  are unchanged.

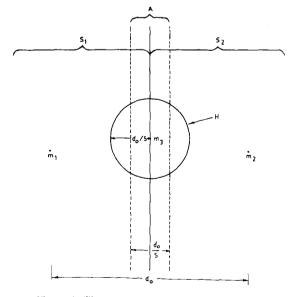


Figure 4. The process of merging in 2-dimension.

Let  $S_{11}$  be such that  $A \cap S_{11} \neq \emptyset$ . It is then generally true that  $A \cap S_{12} = \emptyset$ . Thus  $S_{11}$  and  $S_2$  are to be merged if the merging criterion is satisfied.

4. If both  $S_1$  and  $S_2$  are further subdivided then the merging technique is to be applied similarly.

#### 3. Proposed algorithm

In the proposed algorithm the number of iterations is represented by J while the number of clusters at the Jth iteration is given by  $K_J$ . Initially

$$J = 0$$
 and  $K_0 = 1$ .

Thus, there exists one cluster initially which is S. The algorithm has the following steps.

- Step 1. Apply type-I splitting on every cluster.
- Step 2. If no cluster is divided in Step 1, go to Step3. Otherwise go to Step 1.
- Step 3. Apply the 2-means algorithm on every cluster. Find out the *almost equidistant point* sets for every cluster. Check the inequality (2) for every one of the *almost equidistant point sets* which do not satisfy the *almost equidistant* restriction; then go to Step 4. Otherwise go to Step 5.
- Step 4. Let  $A_1, A_2, ..., A_L$  be those almost equidistant point sets which do not satisfy the almost equidistant restriction. Let the corresponding clusters for  $A_1, A_2, ..., A_L$  be  $C_1, C_2, ..., C_L$ . Let every  $C_s$  be divided into  $C_{s1}$  and  $C_{s2}$  in Step 3 for s = 1, 2, ..., L. Apply type-I splitting on  $C_{s1}$ and  $C_{s2}$  for s = 1, 2, ..., L. For every s = 1, 2, ..., L, the following two cases arise.
  - (a) None of  $C_{s1}$  and  $C_{s2}$  is split. Then check the *merging restriction* on  $C_{s1}$  and  $C_{s2}$  using  $A_{s}$ .
  - (b) At least one of the  $C_{s1}$  and  $C_{s2}$  is split. Then check the *merging restriction* for those sets with which  $A_s$  has non-empty intersection.

Step 5.  $J \leftarrow J+1$ . Find the number of cluster  $K_J$ . Step 6. If  $K_J = K_{J-1}$  then go to Step 7. Otherwise go to Step 1.

Step 7. Stop.

The flowchart of the above algorithm is shown in Figure 5.

Note that the algorithm converges when the

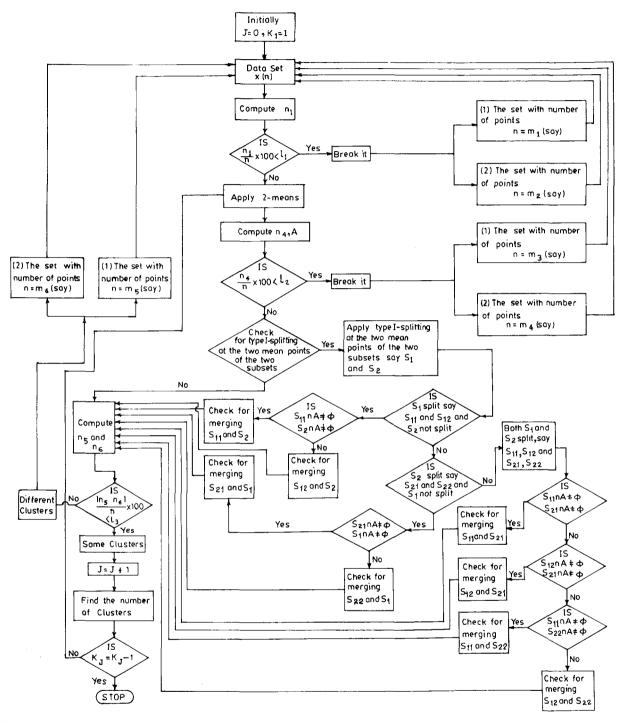


Figure 5. A flowchart of the proposed algorithm.

number of clusters in the Jth iteration and the (J+1)th iteration are equal. A theorem is stated below which shows that if the number of clusters in the Jth and (J+1)th iterations are equal then the

clusters at the end of Jth iteration are the same as the clusters at the end of the (J+1)th iteration. Theorem 1 and Theorem 2 stated below are proved in the Appendix.

**Theorem 1.** Let the number of clusters at the end of the J th and (J+1)th iterations be k and l respectively. Then  $k \leq l$ .

**Theorem 2.** Let the clusters at the end of the Jth and (J+1)th iterations be  $P_1, P_2, ..., P_k$  and  $Q_1, Q_2, ..., Q_k$ . Then

$$\{P_1, P_2, \dots, P_k\} = \{Q_1, Q_2, \dots, Q_k\}.$$

#### 4. Experimental results

#### 4.1. On artificial data

The clustering scheme described in the paper has been implemented on various data sets. The programs are run on an IBM PC/AT microcomputer in TURBO PASCAL language. In all the tested cases, the breadth of the strip, i.e.,  $2h_n = 2a\varepsilon_n$  $[\varepsilon_n = 1/n^p, 0 and <math>a = 3$ , p = 0.05 and n = number of points], the values of  $l_1$ ,  $l_2$  and  $l_3$ are taken 4, 5, 3 and 0.7, respectively.

Figure 6(a) shows multi-cluster data of size 204. For type-I splitting of this data, the mean point is marked by 'X', the direction is marked by a dotted line (3 ---- 7) shown in Figure 6(b). Here  $n_1 = 21$  and also St\_[1 \leftrightarrow 5] and St\_[3 \leftrightarrow 7] contain 21 points each. But

$$\frac{n_1}{n} \times 100 \approx 10.3 > 5$$

so the strip is *densely populated* and is not to be split. The results of type-II splitting are shown in Figure 6(c). The mean points of the two subsets  $S_1$  and  $S_2$  (obtained by using the 2-means algorithm) are marked by ' $X_1$ ' and ' $X_2$ ' and the boundary line is marked by ' $B_1B_2$ '. The *almost equidistant points* are marked by ' $\odot$ '. Here  $n_4=15$ . The mean point of the *almost equidistant points* is marked by 'C'.

Since  $(n_4/n) \times 100 = 7.3 > 3$ , the merging criterion is to be checked. For  $S_1$  and  $S_2$ , the type-I splitting technique is also to be verified. Thus for  $S_1$ , 4 strips are generated with width 4 at the center  $X_1$  and the minimum size of the strip is found for strip St  $[2 \leftrightarrow 6]$  with size 9. Since  $(9/n) \times 100 = 4.4 < 5$ ,  $S_1$  should indeed be split at

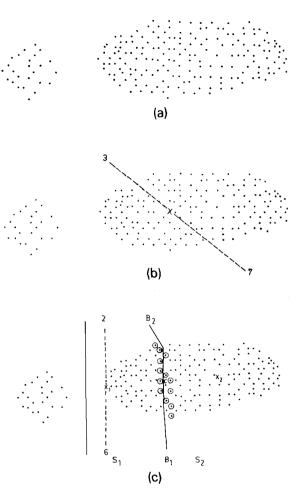


Figure 6. Synthetic data to show the need of type-II splitting. (a) A multi-cluster data of size 204. (b) The mean point and the direction of the low-density strip. (c) The resulting clusters.

the center along the direction  $2 \leftrightarrow 6$ . Similarly it can be seen that  $S_2$  should not be split at the center  $X_2$ .

The equidistant point set, H, for the pair  $S_1$  and  $S_2$  is found. The number of points within the *merging circle* is found to be 25, the values for  $d_0$ ,  $n_5$  and  $n_6$  are found to be 22.3, 13 and 12, respectively. Here

$$\frac{|n_5 - n_6|}{n} \times 100 = 0.4 < 0.7.$$

So one of the divided portions of  $S_1$  is to be merged with  $S_2$ . Thus at the end of the first iteration, the number of clusters is found to be 2 (Figure 6(c)).

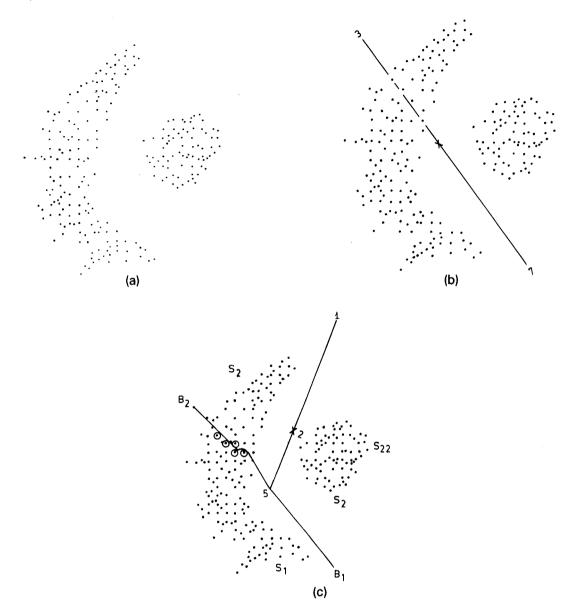


Figure 7. Another example. (a) A multi-cluster data of size 237. (b) The mean point and the direction of the low-density strip. (c) The resulting clusters.

On each one of the above two clusters, type-I splitting and type-II splitting are applied. It has been found that no further divisions can take place. Thus at the end of the second iteration too, the number of clusters is the same. So the process is terminated. One cluster has 25 points and the other has 179 points.

The results of the algorithm on other data sets are also demonstrated. Figure 7(a) shows a data set

and the corresponding intermediary step is shown in Figure 7(b) where the mean point is marked by 'X', the direction is marked by a dotted line (3----7). But the strip along (3----7) is *densely populated* and is not to be split. The results of type-II splitting are shown in Figure 7(c). The boundary line of the two subsets  $S_1$  and  $S_2$  is marked by ' $B_1B_2$ '. For  $S_2$ , the minimum size of the strip is found for strip St\_[1  $\leftrightarrow$  5] in the direction  $1 \leftrightarrow 5$  at the center 'X<sub>2</sub>'. The minimum number of points in St\_[1 \leftrightarrow 5] satisfies the *splitting restriction*. S<sub>2</sub> should indeed be split at the center 'X<sub>2</sub>' along the direction  $1 \leftrightarrow 5$ . S<sub>21</sub> and S<sub>1</sub> are merged. Hence the two clusters are S<sub>22</sub> and S<sub>21</sub>  $\cup$  S<sub>1</sub> which are shown in Figure 7(c).

Figure 8(a) shows another data set which can be split by type-I splitting only and the coresponding results are shown in Figure 8(b). No merging is needed in this case.

The results of applying the proposed algorithm on a real-life data are given below.

#### 4.2. On remotely sensed data

Analysis of satellite imager has wide applications such as crop yield estimation, estimation of

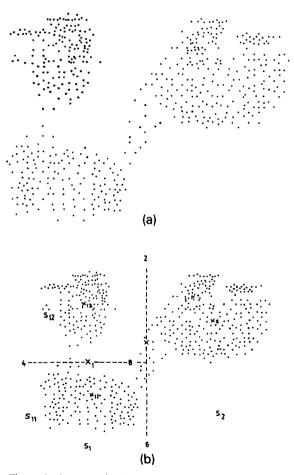


Figure 8. An example where type-I splitting only is needed. (a) The data set. (b) Resulting clusters.

Band	Wavelength	
Blue	0.45 μm-0.52 μm	
Green	$0.52 \mu m$ - $0.59 \mu m$	
Red	0.62 μm=0.68 μm	
Infrared	0.77 μm=0.86 μm	

forest regions etc. [10,11]. Satellite images are used for defence applications too. The proposed algorithm was applied on an Indian remote sensing satellite (IRS) image.

The IRS provides images of two ground resolutions, namely  $72.5 \text{ m} \times 72.5 \text{ m}$  and  $32.5 \text{ m} \times 32.5$ m. It has four bands namely blue, green, red and infrared. The wavelengths of these bands are given in Table 1 [12].

The proposed algorithm has been applied on an IRS image of ground resolution  $36.25 \text{ m} \times 36.25$  m. The area under consideration in that image is a suburb near Calcutta namely Barrackpore. The images corresponding to green and infrared bands for that scene are given in Figures 9 and 10, respectively. The observed gray value in the infrared band has the range 16 to 66 and the green band has the range 18 to 51. The bivariate frequency table of the gray values of these two bands is given in Table 2. The clustering partition by the proposed algorithm is shown as the dashed line in Table 2. The clusters mapped back in the image space are shown in Figure 11. Here, the white pixels in

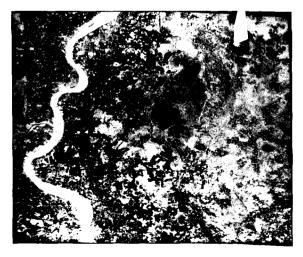


Figure 9. Remotely sensed image corresponding to green band.

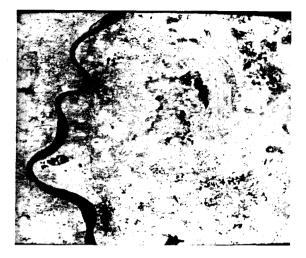


Figure 10. Remotely sensed image corresponding to infrared band.

Figure 11 give the water pixels in the scene. The Hooghly river in the scene is demarkated distinctly from the rest. The path of the river is clearly seen in Figure 11.

#### Acknowledgement

The authors wish to thank Prof. D. Dutta Majumder for his interest in the work. The authors also acknowledge Mr. J. Gupta for typing the manuscript and Mr. S. Chakraborti for his drawings.

### Appendix

**Result 1** [13]. Let  $\varepsilon_n \to 0$  and  $n\varepsilon_n^q \to \infty$ ,  $\varepsilon_n > 0 \quad \forall n$  and q is a positive integer  $\ge 2$ .

Let  $X_1, X_2, ..., X_n$  be independent and identically distributed random vectors following uniform distribution on  $\alpha$ , where  $\alpha \subseteq \mathbb{R}^q$ ,  $\alpha$  is unknown and  $\lambda(\delta \alpha) = 0$  [ $\lambda$  is the Lebesgue measure in q dimensions and  $\delta \alpha$  is the boundary of  $\alpha$ ]. Let

$$\alpha_n = \bigcup_{i=1}^n \{ x \in \mathbb{R}^q \colon \|x - x_i\| \leq \varepsilon_n \}.$$

Then  $\alpha_n$  is a consistent estimator of  $\alpha$ , i.e.,

$$E_{\alpha}[\lambda(\alpha_n \bigtriangleup \alpha)] \to 0 \text{ as } n \to \infty.$$



Figure 11. Resulting two clusters mapped in image domain.

[*E* represents expectation and  $\triangle$  represents symmetric difference.]

If this estimation procedure for a suitably selected  $\varepsilon_n$  is applied to the data in Figure 2(a), then the approximate output is shown in Figure 2(c).

Let  $S \subseteq \mathbb{R}^q$  and  $G = \{x_1, x_2, \dots, x_n\} \subseteq S$ . Now let

$$d_x = \inf_{\substack{y \in G \\ y \neq y}} d(x, y) \quad \forall x$$

where d(x, y) is the distance between two points x and y. Let

$$b = \max_{x \in G} d_x.$$

Then note that if  $b > 2\varepsilon_n$  then the estimated set of G will have at least two components. Thus, G will be partitioned.

Hence, if the data is to be checked for splitting at  $X_0$  of Figure 2(a) then note that  $h_n$  should be greater than or equal to  $\varepsilon_n$ . The choice of  $h_n$  will be  $h_n = a\varepsilon_n$  where  $\varepsilon_n = 1/n^p$ , 0 and a is aconstant.

**Proof of Theorem 1.** Let the clusters at the end of the *J*th and (J+1)th iterations be  $P_1, P_2, ..., P_k$  and  $Q_1, Q_2, ..., Q_l$  respectively.

It is known that there are k clusters in the beginning of the (J + 1)th iteration. Each iteration treats the clusters present at the beginning independently.

Table 2

Bivariate frequency table of gray levels for green and infrared band images

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No part of one cluster becomes a part of another cluster when the iteration ends, i.e.,  $T_1 \subseteq P_{j_1}$ ,  $T_2 \subseteq P_{j_2}$ ,  $j_1 \neq j_2$  and  $T_1 \cup T_2 \subseteq Q_{j_3}$  is not possible  $\forall j_1, j_2$  and  $j_3$ . At most some clusters among  $P_1, P_2, \dots, P_k$  may be split and merged among themselves as the iteration progresses. Thus  $k \leq l$ .  $\Box$ 

**Proof of Theorem 2.** Observe that in the (J+1)th iteration no element of  $P_r$  can change its membership from  $P_r$  to  $P_j$  for  $j \neq r$ , r = 1, 2, ..., k. At most  $P_r$  can be subdivided for r = 1, 2, ..., k.

Let  $P_r$  be split into  $m_r$  subclusters,  $m_r \ge 1$ . If  $m_r = 1$  then there is no split in  $P_r$ . That means, for  $P_r$ , let  $Q_{r_1}, Q_{r_2}, \dots, Q_{r_m}$  be such that

$$\bigcup_{j=1}^{m_r} Q_{r_j} = P_r, \quad r = 1, 2, \dots, k,$$
  
$$m_r \ge 1 \quad \text{and} \quad \sum_{r=1}^k m_r = l.$$

But we know that  $\sum_{r=1}^{k} m_r = k$ . If  $\exists \alpha, \alpha \in \{1, 2, ..., k\}$  such that  $m_{\alpha} > 1$  then  $\sum_{r=1}^{k} m_r > k$ , which is a contradiction. So

$$m_r \leqslant 1 \quad \forall r = 1, 2, \dots, k.$$

But  $m_r \ge 1$   $\forall r = 1, 2, ..., k$  (as stated earlier). Therefore  $m_r = 1, r = 1, 2, ..., k$ . Hence there is no split in  $P_r$ 's. Thus

$$\{P_1, P_2, \ldots, P_k\} = \{Q_1, Q_2, \ldots, Q_k\},\$$

i.e.,  $P_r$  is equal to one of  $Q_1, Q_2, ..., Q_k$   $\forall r = 1, 2, ..., k$ . So the clusters at the end of the *J*th iteration and the (J+1)th iteration are the same.  $\Box$ 

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